A Case for Page Table Reclaim

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Virtual Memory: The pillar of memory management

- 1. Isolation and memory protection between processes
- 2. Gives super powers to the operating system kernel
 - 2.1. Allocating memory on demand (a.k.a demand paging)
 - 2.2. Swapping (to manage overcommitment)
 - 2.3. Transparently mapping memory from different devices (DRAM, CUDA/GPU, NVRAM)
 - 2.4. Copy-On-Write (CoW) for fork
 - 2.5. Various advanced memory management features
 - 2.5.1. NUMA balancing
 - 2.5.2. Huge pages
 - 2.5.3. Virtualisation
 - 2.5.4. Page cache

"All problems in computer science can be solved by another level of indirection." - David Wheeler









Page Table Math

- PTEs can map up to 2MiB of memory
- Each PTE is 4 KiB
- 64 GiB working set requires (64G / 2M * 4K) = 128 MiB
- In general, for X GiB, 2X MiB is required

The Problem of Sparse Mappings

- Libraries like Jemalloc/Tcmalloc leverage sparse mappings for various features
 - Fragmentation avoidance
 - Huge-page aware allocation
- Create a lot of virtual maps over the course of execution
- Call *madvise* (*MADV_FREE*) to free the physical pages, but not the maps
 - To avoid holding mmap locks
 - Avoid rebuilding VMA and page tables
- Page tables accumulate over time and stay unreclaimed

The Problem of Sparse Mappings

- Redis uses jemalloc by default
- On a long-running 600GiB instance of Redis, 100 GiB of physical memory was page tables

The Problem of Memory Hotplug

- Memory desagregation (physical memory allocation on demand) relies on hotplug
- Hotplug is the process of adding/removing memory from a running system
- Memory ballooning also extensively relies on that
- Page tables are unmovable allocations
 - Network buffers
 - Kernel objects
 - Page tables
- Hotunplug fails with unmovable allocations

The Problem of TLB Stalls

- Linux opportunistically uses huge pages with THP (Transparent Huge Page)
 - Reduce TLB stalls
 - Improve access latency
- THP requires allocating higher-order pages (e.g. 2 MiB)
- Unmovable allocations hinders compaction
- Page tables can increase TLB stalls

Reclaiming Page Tables in Linux

- Page tables with no valid references
- Unutilised page tables under memory pressure
 - Host native page tables
 - Guest second-level page tables

Unreferenced Page Tables

- Add a per-cpu reference counter to the page structure
- Whenever an entry is mapped, the ref counter is atomically incremented
- Whenever an entry is unmapped, the refconter is atomically decremented
- Once it reaches zero, the refcount is freezed
 - Concurrent threads won't get a ref
- PMD entry is cleared and page table is freed

PTE Reclaim under Memory Pressure

- Second-level PTEs for VMs always contain valid entries
- If the system is short on memory, page tables can be reclaimed
- All PTEs are inserted into a per-memcgroup list upon allocation
- Memory pressure randomly picks PTEs from the list
 - Freeze the refcount (to prevent concurrent reclaim)
 - Empty and zero page references only are freed
 - PTEs with valid references are
 - Migrated to another NUMA node
 - Compressed in-memory
 - Swapped out to disk

Observations

- On a microbenchmark with madvise(), went from 100 MiB PTEs to 100 KiB
- No overhead for refcount reclaim
- Swapping helps manage memory overcommitment situations
 - Reduced OOM killer invocations

Ongoing/Future work

- It is only applied to anonymous mapping for now
 - Page cache on the way
- PTE list is not LRU sorted
 - Need a mechanism to track access frequency of page tables
- Leverage migration for compaction
- How much does it help for reducing TLB stalls?



Thanks! Questions?